

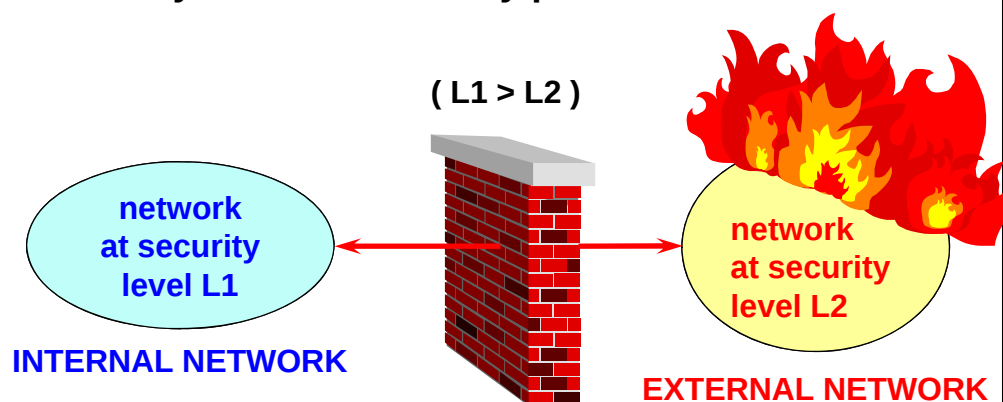
Firewall and IDS/IPS

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What is a firewall?

- firewall = wall to protect against fire propagation
- controlled connection between networks at different security levels = boundary protection



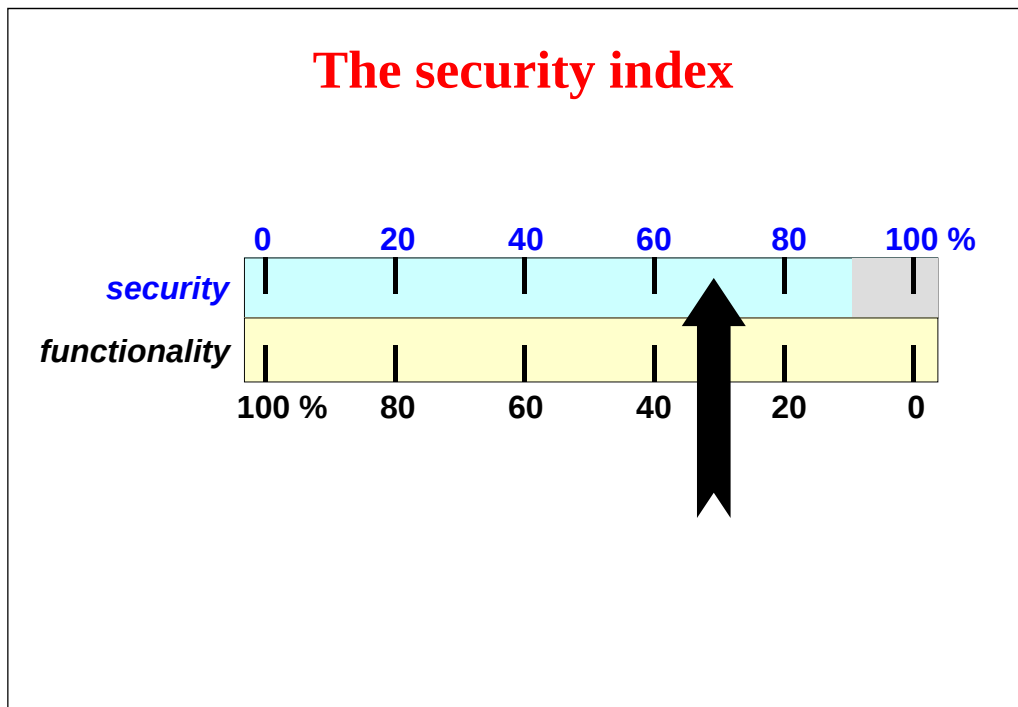
Ingress vs. Egress firewall

- **ingress firewall**
 - incoming connections
 - typically to protect my public services
 - sometimes as part of an application exchange initiated by my users
- **egress firewall**
 - outgoing connections
 - typically to check the activity of my personnel (!)
- **easy classification for channel-based services (e.g. TCP applications), but difficult for message-based services (e.g. ICMP, UDP applications)**

Firewall design

**You don't "buy" a firewall, you design it!
(you can buy its components)**

- **we need to achieve an optimal trade-off ...**
- **... between security and functionality**
- **... with minimum cost**



THE THREE COMMANDMENTS OF FIREWALL

- I. the FW must be the only contact point of the internal network with the external one
- II. only the “authorized” traffic can traverse the FW
- III. the FW must be a highly secure system itself

*D.Cheswick
S.Bellovin*

Authorization policies

**“All that is not
explicitly permitted,
is forbidden”**

- higher security
- more difficult to manage

**“All that is not
explicitly forbidden,
is permitted”**

- lower security (open gates)
- more easy to manage

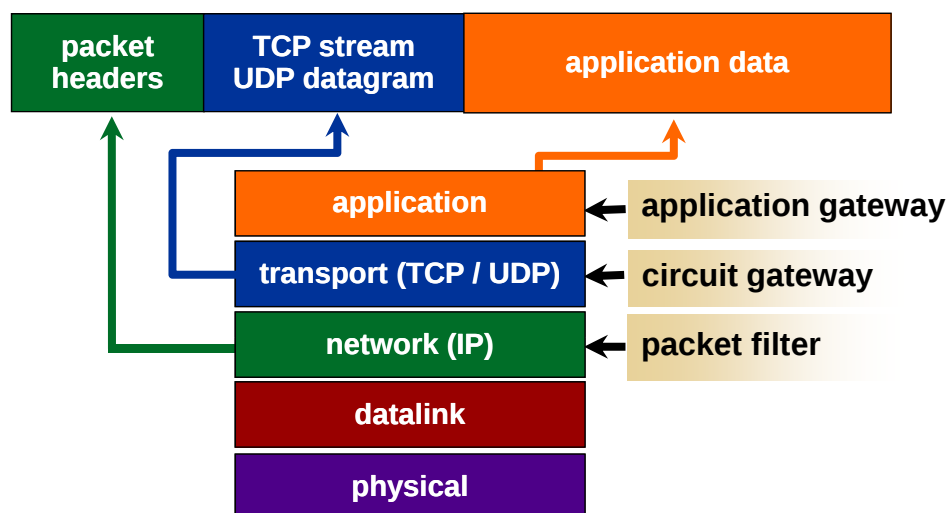
General concepts

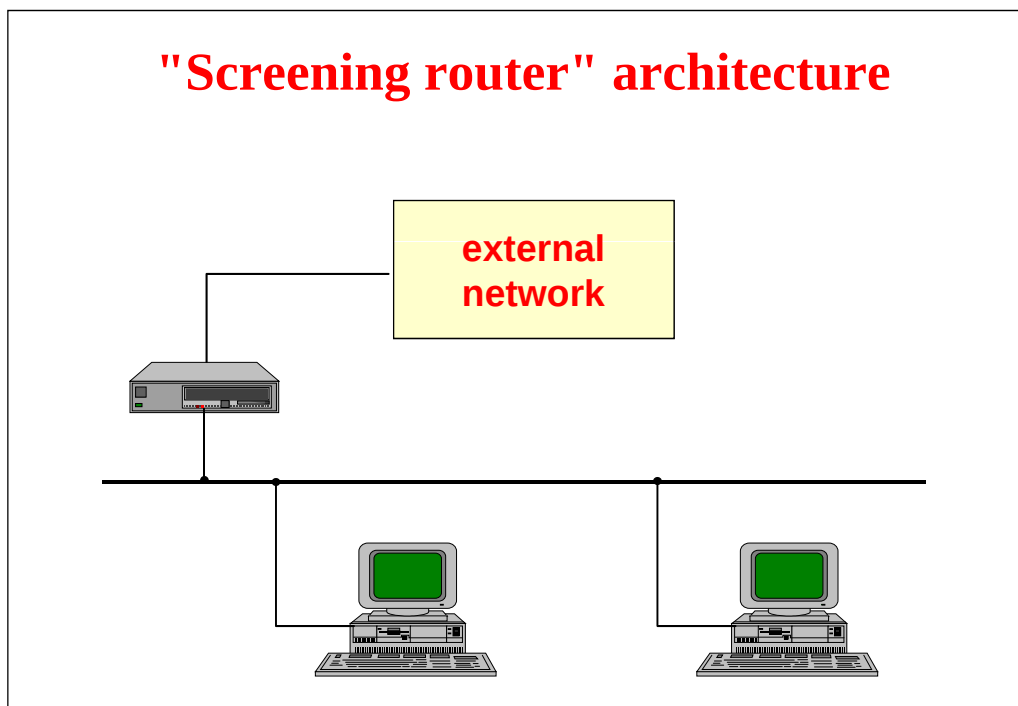
- the bigger an object, the more difficult to verify it
- if a process has not been activated, its bugs are not relevant
- “big is NOT beautiful” = minimal configuration
- a FW is not a general-purpose machine:
 - minimal sw
 - no users
 - ...
- everybody is guilty unless he proves his innocence

FW: basic components

- **screening router (choke)**
router that filters traffic at IP level
- **bastion host**
secure system, with auditing
- **application gateway (proxy)**
service that works on behalf of an application, with access control
- **dual-homed gateway**
system with two network cards and routing disabled

A which level the controls are made?

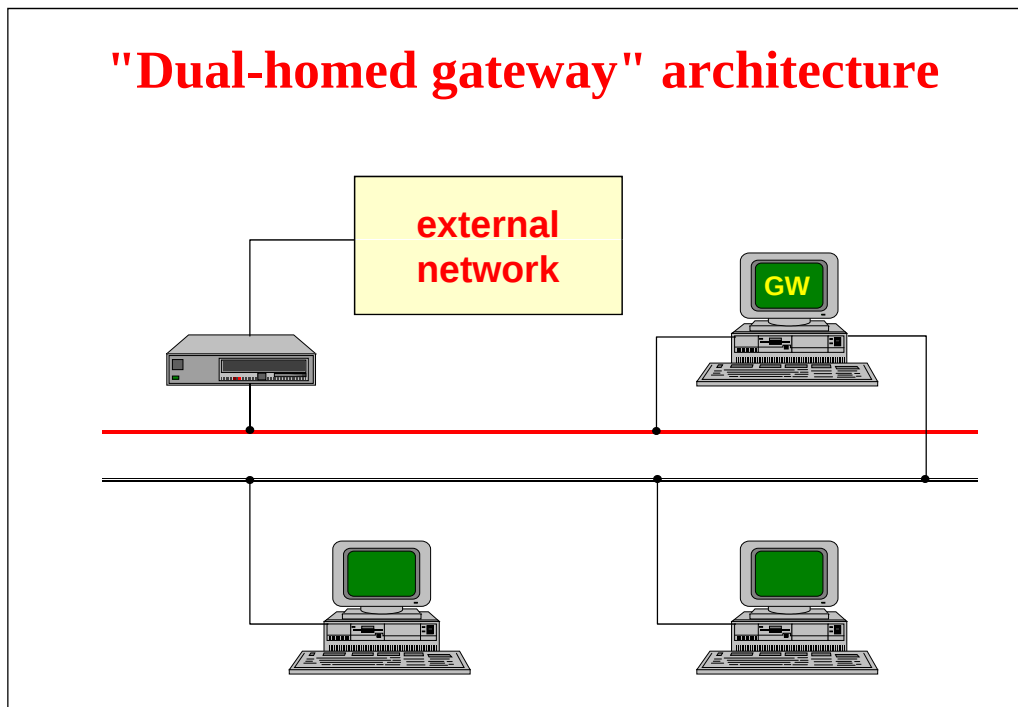




"Screening router" architecture

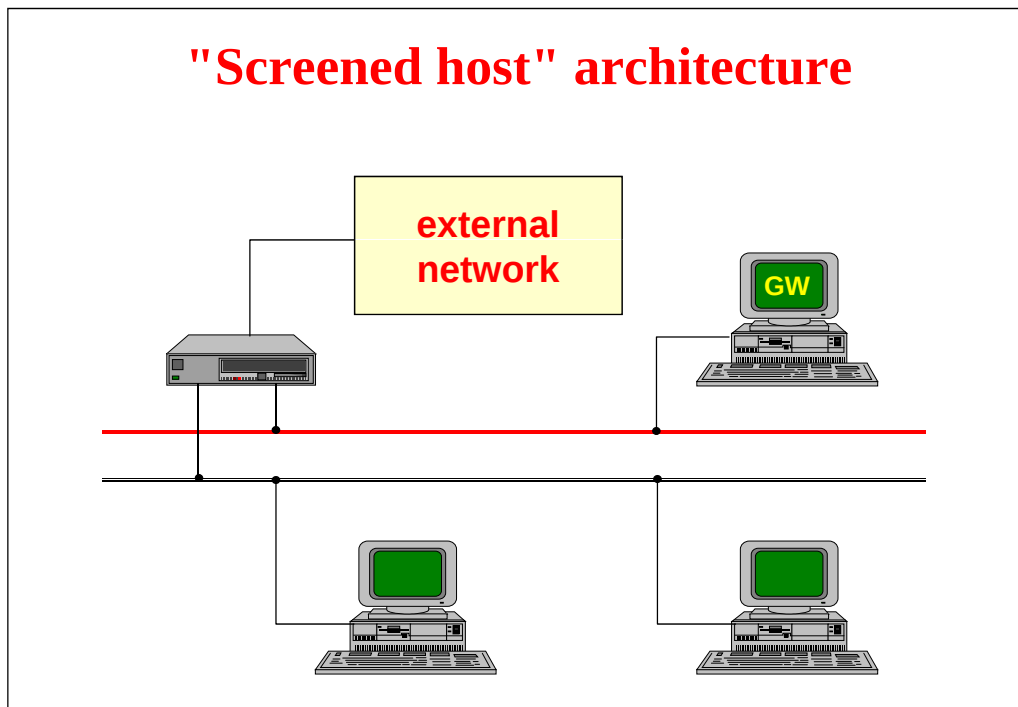
- exploits the router to filter the traffic both at IP and upper levels
- no need for dedicated hardware
- no need for a proxy and hence no need to modify the applications
- simple, easy, cheap and ... insecure!

"Dual-homed gateway" architecture



"Dual-homed gateway" architecture

- easy to implement
- small additional hardware requirements
- the internal network can be masqueraded
- unflexible
- large work overhead



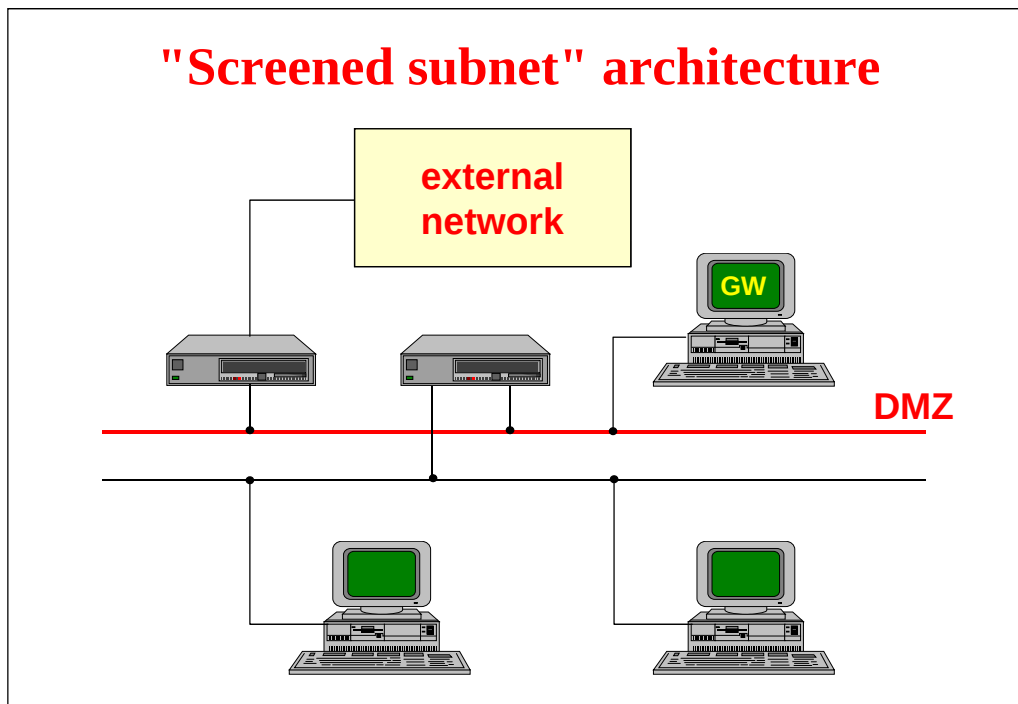
"Screened-host" architecture

- **router:**
 - blocks the packets from INT to EXT unless they come from the bastion host
 - blocks the packets from EXT to INT unless they go to the bastion host
 - exception: directly enabled protocols
- **bastion host:**
 - circuit/application level gateway to selectively enable some services servizi

"Screened-host" architecture

- more expensive
- more flexible
- complex to manage: two systems rather one
- possible selectively relax the controls over some services / hosts
- only the hosts/protocols passing through the bastion can be masqueraded (unless the router offers the NAT functionality)

"Screened subnet" architecture

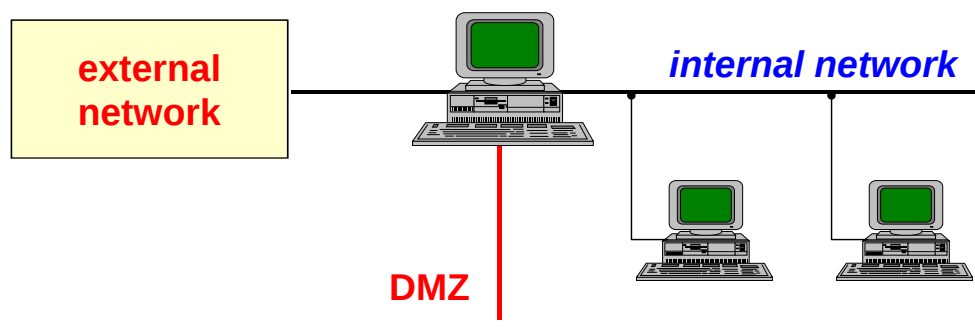


"Screened subnet" architecture

- DMZ (De-Militarized Zone)
- the DMZ is home not only to the gateway but also to other hosts (typically the public servers):
 - Web
 - remote access
 - . . .
- the routing may be configured so that the internal network is unknown
- expensive

"Screened subnet" architecture (version 2)

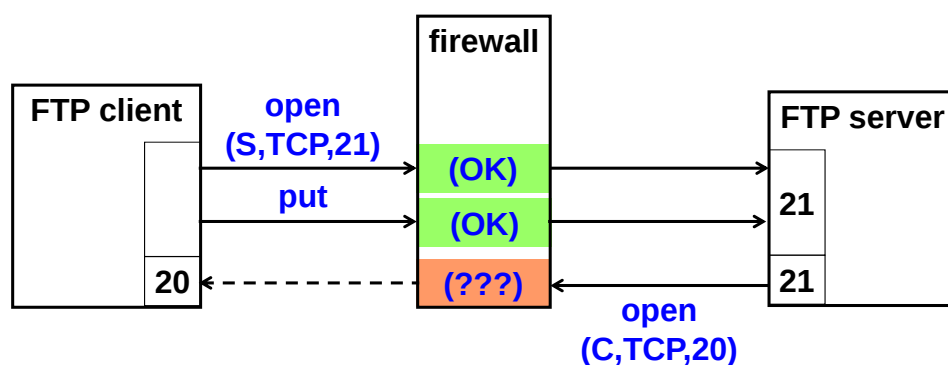
- to reduce costs and simplify management often the routers are omitted (and their function incorporated into the gateway)
- AKA "three-legged firewall"

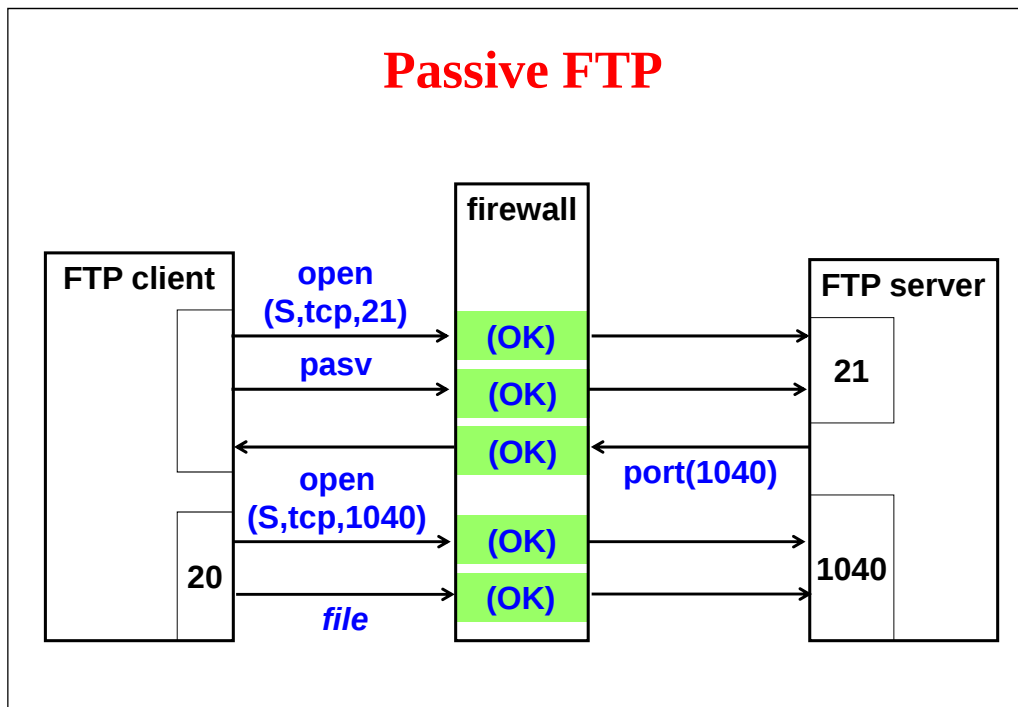


Filters at network level (I)

- address checking (ingress / egress filtering)
- disable incoming services (with exceptions):
 - e.g. TELNET only towards INET
 - e.g. incoming HTTP only to the DMZ web server
 - problem with FTP (the data transfer channel is always created by the server)
- ICMP
 - is dangerous (used for denial-of-service) but useful (ping, traceroute) so don't disable it, just rate-limit it
 - closely monitor REDIRECT packets

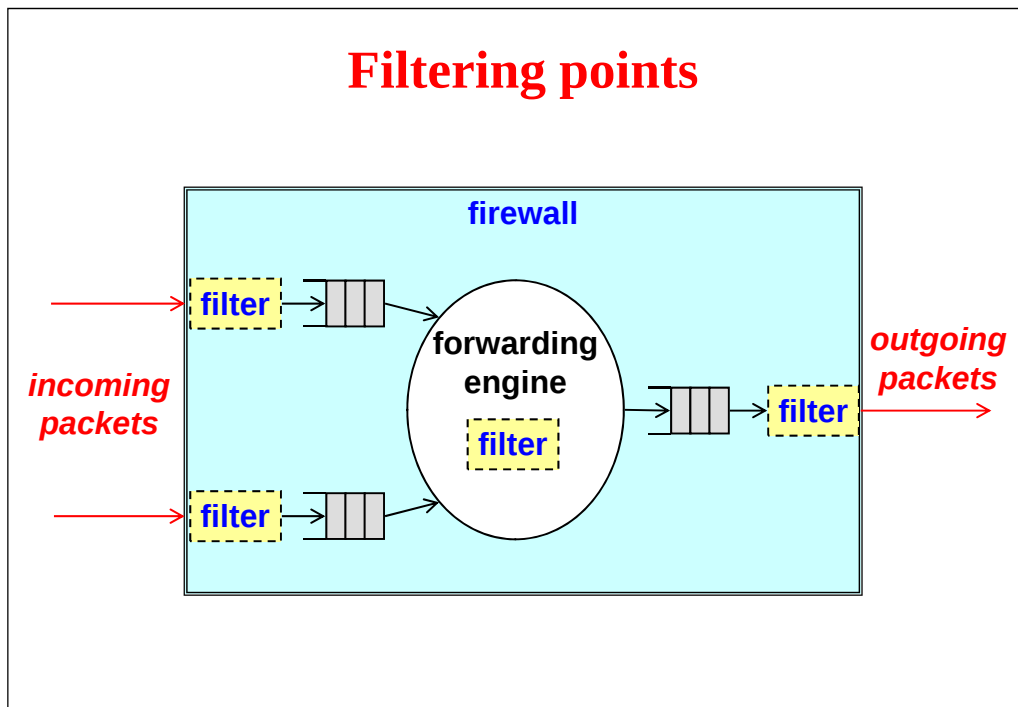
Problem with FTP across a firewall





Filters at network level (II)

- **UDP**
 - is a datagram service, not a virtual circuit (so its checking poses a higher load)
 - RPC uses random ports
 - it is suggested to completely disable it (but DNS)
- **distinguish internal and external interfaces**
- **pay attention to the number of filtering rules and their order: this may drastically change performance**



Filters on a router: an example

- **policy:** mail for network 130.193 managed only by 130.193.2.1
- **syntax of CISCO router:**
 - access-list 100 permit tcp
0.0.0.0 255.255.255.255
130.193.2.1 0.0.0.0
eq 25
 - access-list 101 deny tcp
0.0.0.0 255.255.255.255
130.193.0.0 0.0.255.255
eq 25

Bastion host - configuration

- should run only the relevant processes
- must log (securely!) all its activities
- network log onto an internal secure system
- must have *source routing* disabled
- must have *IP forwarding* disabled
- should have “mouse traps” (e.g. fake 1s)

Firewall technologies

- **different controls at various network levels:**
 - (static) packet filter
 - stateful (dynamic) packet filter
 - cutoff proxy
 - circuit-level gateway / proxy
 - application-level gateway / proxy
 - stateful inspection
- **differences in terms of:**
 - performance
 - protection of the firewall O.S.
 - keeping or breaking the client-server model

Packet filter

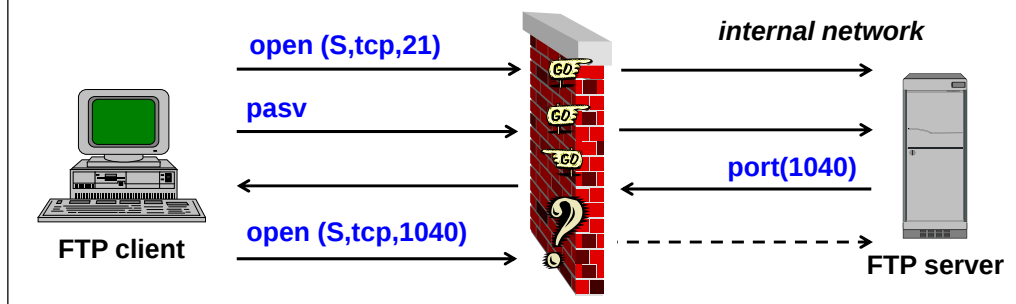
- **historically available on routers**
- **packet inspection at network level**
 - IP header
 - transport header

Packet filter: pros and cons

- **independent of applications**
 - good scalability
 - approximate controls: easy to “fool”
(e.g. IP spoofing, fragmented packets)
- **good performance**
- **low cost (available on routers and in many OS)**
- **difficult to support services with dynamically allocated ports (e.g. FTP)**
- **complex to configure**

Packet filter & FTP

- **two choices available:**
 - leave open all dynamic ports (>1024)
 - close all dynamic ports
- **difficult trade-off between security and support to FTP!!**



Stateful (dynamic) packet filter

- **similar to packet filter but “state-aware”**
 - state informations from the transport or application level (e.g. FTP PORT command)
 - can distinguish new connections from those already open
 - state tables for open connections
 - packets matching one row in the table are passed without any further control
- **better performance than packet filter**
 - SMP support
- **still has many of the static packet filter limitations**

Application-level gateway

- composed by a set of proxies inspecting the packet payload at application level
- often requires modifications to the client application
- may optionally mask / renumber the internal IP addresses
- when used as part of a firewall, usually performs also peer authentication
- top security!! (e.g. against buffer overflow of the target application)

Application-level gateway (1)

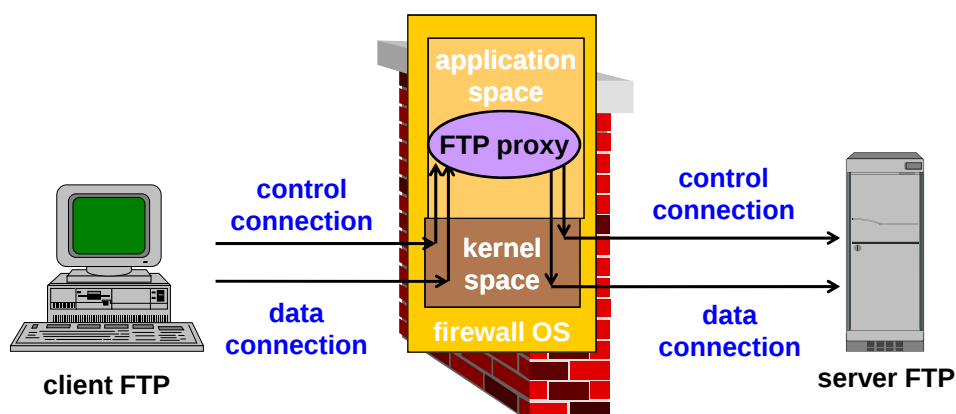
- rules are more fine-grained and simple than those of a packet filter
- every application needs a specific proxy
 - delay in supporting new applications
 - heavy on resources (many processes)
 - low performance (user-mode processes)
- SMP may improve performance
- completely breaks the client/server model
 - more protection for the server
 - may authenticate the client
 - not transparent to the client

Application-level gateway (2)

- the fw OS may be exposed to attacks
- problems with application-level security techniques (e.g. SSL)
- variants:
 - transparent proxy
 - less intrusive for the client
 - strong application proxy
 - only some commands/data are forwarded
 - this is the only correct configuration for an application proxy

Application-level gateway & FTP

- total control of the application session



Circuit-level gateway (1)

- **a generic proxy (i.e. not “application-aware”)**
 - creates a transport-level circuit between client and server ...
 - ... but it doesn't understand or manipulate in any way the payload data

Circuit-level gateway (2)

- **breaks the TCP/UDP-level client/server model during the connection**
 - more protection for the server
 - isolated from all attacks related to the TCP handshake
 - isolated from all attacks related to the IP fragmentation
 - may authenticate the client
 - but this requires modification to the application
- **still exhibits many limitations of the packet filter**

SOCKS

- a transport-level proxy (L4) that implements a circuit-level gateway
- invented at MIPS, v4 by NEC, v5 by IETF
- aka AFT (Authenticated Firewall Traversal)
- requires modified clients:
 - standard: telnet, ftp, finger, whois
 - library to develop own clients
- commercial support:
 - all the major browsers (e.g. FX and IE)
 - some firewalls (e.g. IBM)

SOCKS RFCs

- RFC-1928 “SOCKS protocol V5”
- RFC-1929 “Username/password authentication for SOCKS V5”
- RFC-1961 “GSS-API authentication method for SOCKS V5”
- RFC-3089 “A SOCKS-based IPv6/IPv4 gateway mechanism”

SOCKS: how it works

- the library replaces the standard socket system calls `connect()`, `bind()`, `accept()`, ...
- ... with calls that:
 - open a channel to the SOCKS server
 - send `version`, `IP:port`, `user`
- the server:
 - checks its ACL
 - opens the required channel (with its own IP address) and relays L4 payloads

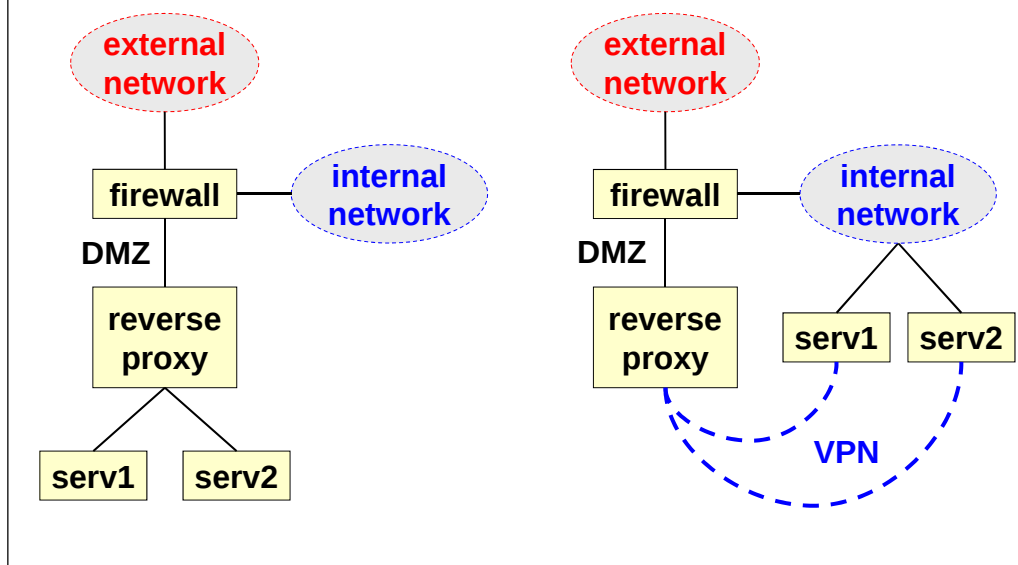
SOCKS: initial problems

- **SOCKS v4:**
 - doesn't distinguish the internal and external nets
 - weak user authentication (based on `identd` or `id` stored locally on client)
 - supports only TCP
- **solution = SOCKS v5:**
 - supports UDP too
 - better authentication (user+pwd or GSS-API)
 - encryption (between the SOCKS client and server)

Reverse proxy

- HTTP server acting just as a front-end for the real server(s) which the requests are passed to
- benefits:
 - obfuscation (no info about the real server)
 - load balancer
 - SSL accelerator (with unprotected back-end connections ...)
 - web accelerator (=cache for static content)
 - compression
 - spoon feeding (gets from the server a whole dynamic page and feeds it to the client according to its speed, so unloading the application server)

Reverse proxy: possible configurations



FW architectures: which is the right one?

- **theoretically, the higher the level and:**
 - the higher the CPU cycles needed
 - the higher the protection level
- **the (sad) reality:**

Firewall customers once had a vote, and voted in favor of transparency, performance and convenience instead of security; nobody should be surprised by the results.
(Marcus J. Ranum, the “grandfather of firewalls”,
firewall wizard mailing list, oct 2000)

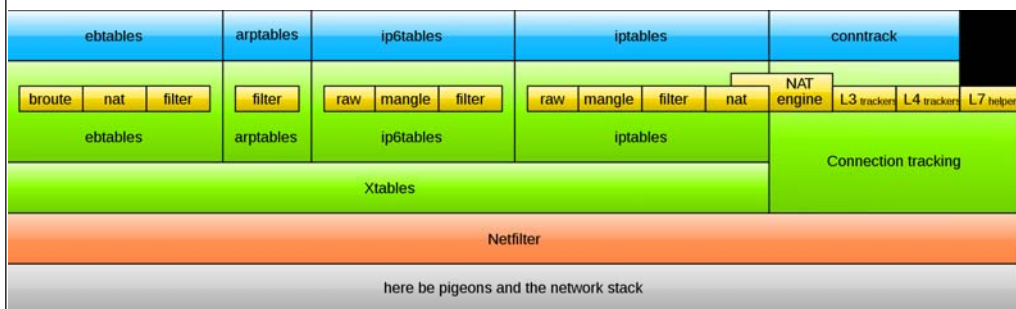
FW architectures: which is the right one?

- **the best choice:**
 - not a single product, but a robust fw architecture that supplements the holes and vulnerabilities of the single components!!!
 - for each component choose something available on several architectures: it's better to be able to choose rather than leaving the choice to the vendor!!
 - beware of solutions promising to solve each and every security problem: may be it's just advertisement ...

Firewall: commercial products

- there is plenty of firewall manufacturers/vendors
- typically on UNIX, sometimes on Windows (the latter requires changing the TCP/IP stack!)
- the free Firewall Toolkit (FWTK)
 - TIS (www.tis.com)
 - base application-gateway components
- the free IPchains / IPfilter / IPtables on Linux
 - packet-filter

Linux: netfilter components



(image from wikipedia)

Default netfilter chains

■ PREROUTING

- all incoming packets, before any routing decision

■ INPUT

- packets with destination the node itself (the "local-delivery" routing table: "ip route show table local")

■ FORWARD

- packets to be forwarded, after the routing decision

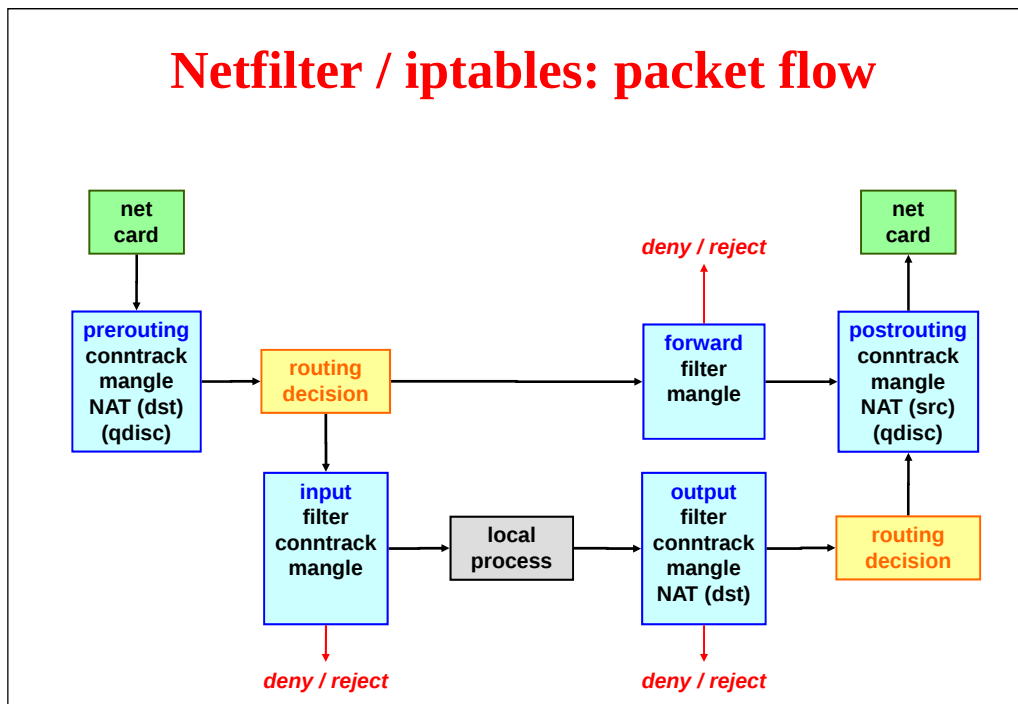
■ OUTPUT

- packets generated by the node itself

■ POSTROUTING

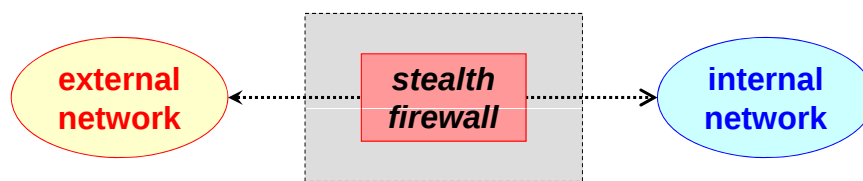
- all outgoing packets, just before sending them

Netfilter / iptables: packet flow



Stealth firewall

- firewall without an IP address, so that it cannot be directly attacked
- physical packet interception (by setting the interfaces in promiscuous mode)

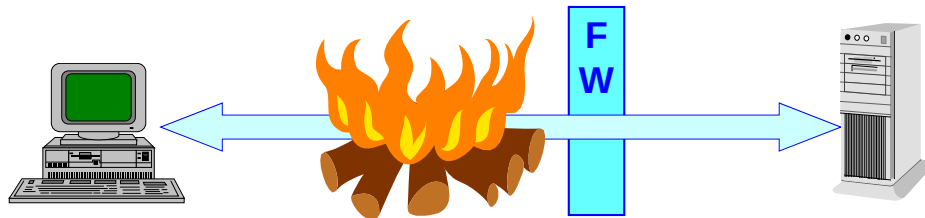


Local / personal firewall

- firewall directly installed at the node to be protected
- typically a packet filter
- w.r.t. a normal network firewall, it may limit the programs than are allowed:
 - to open network channels towards other nodes (i.e. act as a client)
 - to answer network requests (i.e. act as a server)
- important to limit the diffusion of malware and trojans, or plain configuration mistakes
- firewall management must be separated from the system management

Protection offered by a firewall

- a firewall is 100% effective only for attacks over/against blocked channels
- the other channels require other protection techniques:
 - VPN
 - “semantic” firewall / IDS
 - application-level security



Intrusion Detection System (IDS)

- **definition:**
 - system to identify individuals using a computer or a network without authorization
 - extendable to identify authorized users violating their privileges
- **hypothesis:**
 - the behavioural “pattern” of non-authorized users differs from that of authorized users

IDS: functional features

- **passive IDS:**
 - cryptographic checksum (e.g. tripwire)
 - pattern matching (“attack signature”)
- **active IDS:**
 - “learning” = statistical analysis of the system behaviour
 - “monitoring” = active statistical info collection of traffic, data, sequences, actions
 - “reaction” = comparison against statistical parameters (reaction when a threshold is exceeded)

IDS: topological features

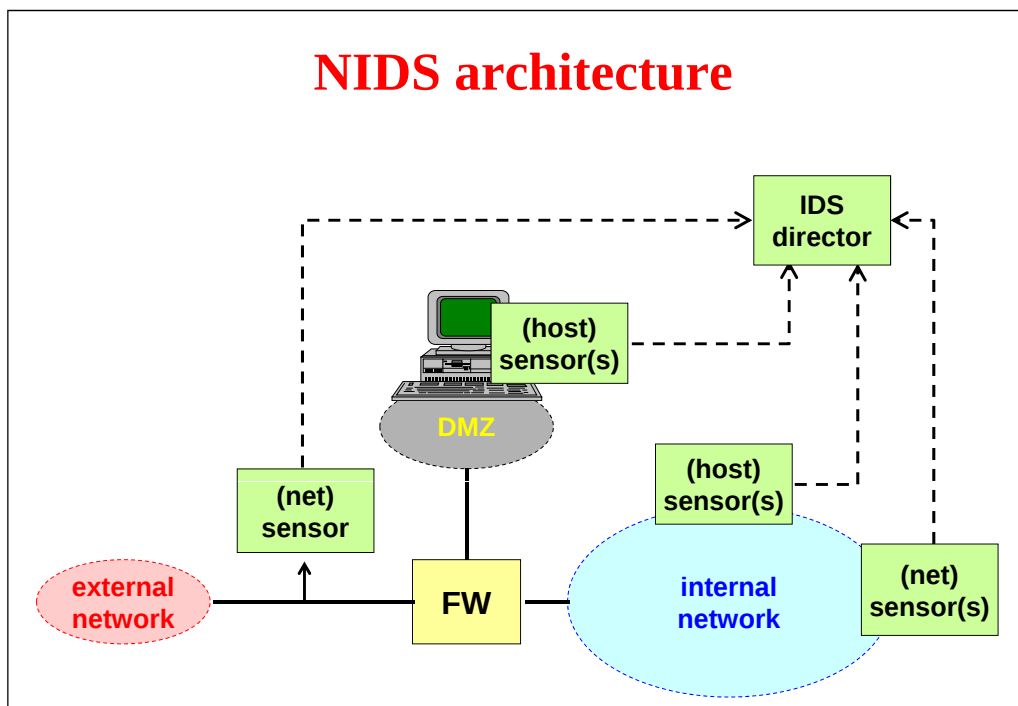
- **HIDS (host-based IDS)**
 - log analysis (OS, service or application)
 - internal OS monitoring tools
- **NIDS (network-based IDS)**
 - network traffic monitoring tools

SIV and LFM

- **System Integrity Verifier**
 - checks files / filesystems looking for changes
 - e.g. changes to Windows registry, cron configuration, user privileges
 - e.g. tripwire
- **Log File Monitor**
 - checks the log files (OS and applications)
 - looks for known patterns of successful attacks or attempts
 - e.g. swatch

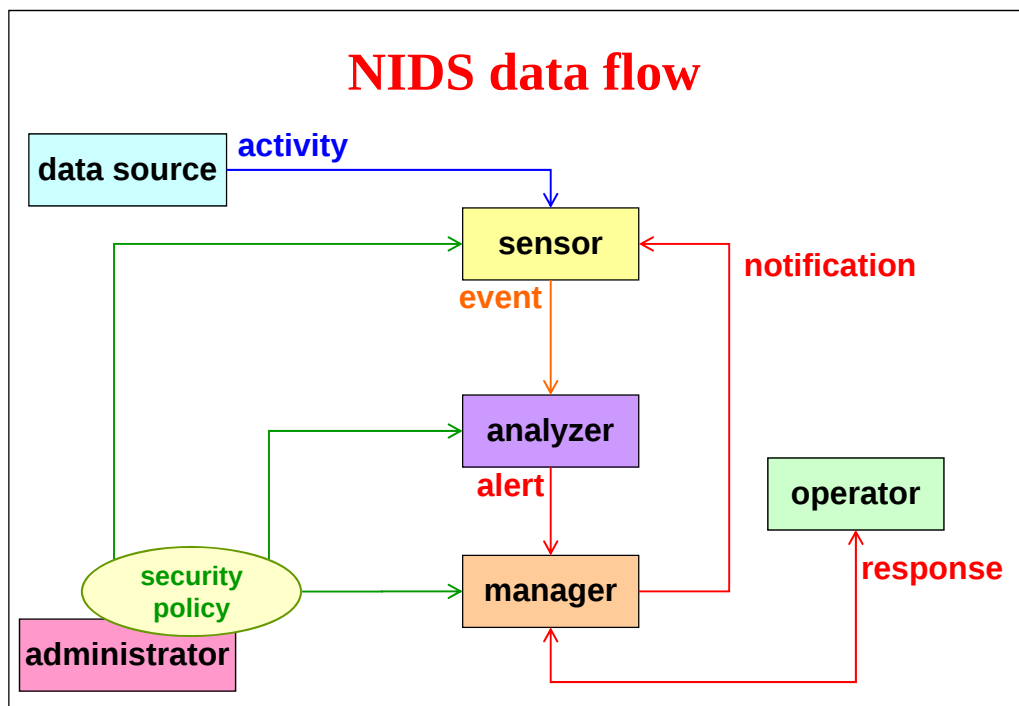
NIDS components

- **sensor**
 - checks traffic and logs looking for suspect patterns
 - generated the relevant security events
 - interacts with the system (ACLs, TCP reset, ...)
- **director**
 - coordinates the sensors
 - manages the security database
- **IDS message system**
 - secure and reliable communication among the IDS components



IDS/NIDS interoperability

- much needed because the attacks involve different organizations and/or are detected by different tools
- attack signature format:
 - no standard, but Snort format is in large use
- alarm format and protocol for alarm transmission:
 - IDMEF + IDXP + IODEF (IETF)
 - SDEE (Cisco, ISS, SourceFire)



IDMEF + IDXP

- developed by the IETF
- **Intrusion Detection Message Exchange Format**
 - independent from the protocol (IPv4, IPv6)
 - supports internationalization and localization
 - supports data aggregation and filtering (on the manager)
- **Intrusion Detection eXchange Protocol**
 - based on BEEP (RFC-3080)
 - profile exchange (for end-to-end security, for ID)
 - base security profile is TLS

SDEE

- **Secure Device Event Exchange**
- **based on the webservice paradigm:**
 - messages in XML
 - messages transported over HTTP or HTTPS
- **closed standard (?), managed by the ICSAlabs**

IODEF

- **Incident Object Description and Exchange Format**
- **is a superset of IDMEF**
- **used to exchange information between different organizations, keep statistics, evaluate risks, ...**

IPS

- **Intrusion Prevention System**
- **to speed-up and automate the reaction to intrusions**
= IDS + distributed dynamic firewall
- **a technology, not a product, with large impact on**
many elements of the protection system
- **dangerous! may take the wrong decision and block**
innocent traffic

